Introduction to High Performance Computing

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SDS406 – Fall semester, 2024 - 2025

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L01: Introduction, 7th October 2024

SDS406

SDS406: Introduction to high performance computing

- Day and time?
- Interleaved lecture and hands-on \Rightarrow please always be with your laptops
- MSc mandatory course
- Elective for PhD students (CyI and UCY)

Assessment

- Coursework
 - Three homework assignments (35%)
 - In-class participation in labs and small homework exercises (15%)
- Final examination (50%)
 - Final assignment

Preliminaries

- Laptop computers for exercises
 - To log in to educational cluster
 - Need an ssh client
 - Need to learn to use a text editor on a remote system
 - e.g., <u>VS Code</u>, <u>Emacs</u>, <u>Vim</u>
 - Go through lesson slides <u>sds406.online</u>
- "Lab" format
 - Each week's lesson will include taught components and practical exercises
 - $\circ~$ For the exercises we will use C (or C++) and Python:
 - C for performance-targeted exercises
 - Python for post-processing (analysis and visualization of results)
- Access to the educational cluster will be provided with dedicated course accounts
 - Accounts prepared for you before this lesson
 - You will use these accounts during the in-class exercises, homeworks, and assignments
- I assume that:
 - this **is not** your first *programming* course
 - this **is** your first *parallel programming* course

This Lesson

Preliminaries on parallel computing

- Terminologies and definitions
- Overview of high performance computing landscape

Preparation for labs

- Log into educational system
- Editing source code files remotely

"The use of supercomputers to solve complex computational tasks"



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(Super)Computing evolution

1940s — First computers, e.g. ENIAC

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1970s, 80s — Dawn of personal computing, but supercomputers needed for specialized tasks, e.g. Cray-1

• parallelism emerges: vectorization, parallel instructions



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1990s, 2000s — Supercomputers via integration of commodity components (e.g. Beowulf clusters)

• Parallelism as we understand it today: distributed and shared memory, message passing, etc.



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• Parallelism as we understand it today: distributed and shared memory, message passing, etc.

2010s onwards — Heterogeneous supercomputers

• Many sockets per node; co-processors, e.g. GPUs, potentially multiple architectures within same system

No Dennard scaling after ~ 2005

- $P \propto A f V^2$ (P: power, A: area, f: freq., V: voltage)
- Roughly, as transistors get smaller, the *power density* stays constant
- This efficiency was typically used towards increasing frequency

https://github.com/karlrupp/microprocessor-trend-data

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What about Moore's law?

- Moore's law is a statement about transistor count
- And so far seems to be holding strong
- \Rightarrow Efficiencies are now used to increase parallelism rather than frequency
 - HPC now means more parallel hardware, rather than faster scalar hardware
 - HPC primarily deals with tackling the challenges of parallelism, from all aspects (hardware, software, algorithmic, etc.)

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Evolution of Supercomputing

The "Top500" list

- Ranked list of top 500 supercomputers populated twice per year
- Data shown here since 1993
- Shown is High-Performance Linpack (HPL) performance achieved
- \Rightarrow Shows that exponential increase has been maintained

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- Shown is High-Performance Linpack (HPL) performance achieved
- \Rightarrow Shows that exponential increase has been maintained
 - Proliferation of accelerators
 - Latest AMD accelerators are Instinct GPUs
 - Intel between 2012 and 2017 includes Xeon Phi
 - IBM accelerators refer to IBM Cell

 \Rightarrow Since 2021, more than half of performance over Top500 now from accelerated systems

https://www.top500.org

Parallel computing as the main paradigm

High performance computing means parallel computing

• Technology trends mean parallelism is essential for advanced computing

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- Practitioners of computational science and engineering benefit from knowledge of concepts and challenges of parallel computing

Parallel computing as the main paradigm

High performance computing means parallel computing

- Technology trends mean parallelism is essential for advanced computing
- Practitioners of computational science and engineering benefit from knowledge of concepts and challenges of parallel computing
 - Architectures and their characteristics
 - Algorithms and how amenable they are to parallelisation
 - Performance metrics and their significance, e.g. sustained and peak floating point performance, bandwidth, scalability

Characterization of Parallel Workloads

Task vs Data Parallelism

Taxonomy of computer architectures, *Flynn's Taxonomy*:

- Single Instruction stream, Single Data stream (SISD)
- Multiple Instruction streams, Single Data stream (MISD)
- Single Instruction stream, Multiple Data streams (SIMD)
- Multiple Instruction streams, Multiple Data streams (MIMD)

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Parallel computing

- Most computing devices have underlying SIMD architecture units: GPUs, CPUs, etc.
- Most supercomputers can be considered MIMD architectures: multiple interconnected computing devices that can issue the same or different instructions on multiple data

Shared vs Distributed memory paradigm

Shared memory

- Multiple processes share common memory (common *memory address space*)
- E.g. multi-core CPU, multi-socket node, GPU threads, etc.
- Programming models: OpenMP, pthreads, MPI, CUDA

Core Core Core Cache Memory

Distributed memory

- Processes have distinct memory domains (different *memory address space*)
- E.g. multiple nodes within a cluster, multiple GPUs within a node
- Programming models: MPI

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Data shared via explicit communication over a network

Data shared via memory

Capacity vs Capability

How do we "spend" parallelism?

- Capacity computing
 - Improve time-to-solution of a problem that can also run on less number of processes
 - E.g. solve many small problems
- Capability computing
 - Solve a problem that was *impossible* to solve on less processes
 - $\circ~$ E.g. solve a problem using N~ nodes, that cannot fit in memory of less nodes

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More "capacity-like"

More "capability-like"

Defining Scalability

Concepts of parallel computing

- Scalability: The rate at which time-to-solution improves as we increase processing units
- Weak scaling: Increase processing units; keep the local problem size fixed \Rightarrow for increasing global problem size
- Strong scaling: Increase processing units; keep the global problem size fixed \Rightarrow for decreasing local problem size

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Quantifying Scalability

Scalability

• Speedup S when using N processes

$$S(N) = \frac{T_0}{T(N)},$$

- T₀: Reference time-to-solution (using N₀ processors, nodes, GPUs, etc.)
- \circ T(N): Time-to-solution using N > N₀ processes
- Parallel efficiency ϵ

$$\varepsilon = S(N) \frac{N_0}{N}$$

• Ideal scaling: $\varepsilon \simeq 1$

Ν	T [hrs]	S	e
28	8.02		
49	4.87	1.65	0.94
56	4.10	1.96	0.98
98	2.67	3.01	0.86
112	2.39	3.36	0.84

Modeling Scalability

Amdahl's Law

- f: fraction of application that can be parallelized
- T₀: time-to-solution of code when using one process
- N: Number of processes

$$T(N) = (1 - f)T_0 + f\frac{T_0}{N}$$
$$S(N) = \frac{T_0}{T(N)} = \frac{1}{1 - f + \frac{f}{N}}$$

T(N)

A practical example

- Consider the calculation of π via simple Monte Carlo:
 - $\circ~$ Define a unit square. Set $n_{hit}=0$
 - \circ Randomly pick points (x, y) within the unit square

$$\circ \ \text{If} \ x^2 + y^2 < 1, n_{hit} + = 1$$

- $\circ \ \ \text{Repeat} \ N \ \text{times}$
- The ratio n_{hit}/N approaches the area of a circle quadrant $\Rightarrow \frac{\pi}{4}$

A practical example

```
unsigned long int hits = 0;
for(unsigned long int i=0; i<N; i++) {
    double x = drand48();
    double y = drand48();
    if(x*x + y*y < 1)
        hits += 1;
}
```

>

A practical example

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• Parallelizable parts

• The loop over N

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• Parallelizable parts

• The loop over N

- Scalar parts
 - Initialization
 - $\circ~$ Summing the partial <code>nhit</code> and division by N

Monte Carlo estimation of π

• N = 14,968,800

n _{proc}	t (sec)	π
1	0.53	3.141543
2	0.27	3.141355
3	0.19	3.141089
4	0.16	3.141137
5	0.12	3.141135
6	0.10	3.141290
7	0.11	3.141533
8	0.11	3.141195

Monte Carlo estimation of π

• N = 479,001,600

n _{proc}	t (sec)	π
1	13.70	3.141558
2	7.58	3.141605
3	4.74	3.141603
4	3.56	3.141647
5	2.86	3.141650
6	2.40	3.141604
7	2.08	3.141625
8	1.82	3.141646

By now you should have followed instructions to:

ightarrow Generate an ssh key-pair

ightarrow Log into our system



- Compute Nodes
 - Computational units CPU and potentially a co-processor, e.g. a GPU
 - Memory (i.e. RAM)
 - Some storage and/or NVMe
 - Network interfaces, possibly separate between management and workload



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 - $\circ \ \ \text{Interfaces on nodes}$
 - Wiring and switches



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- Storage
 - Still predominantly spinning disks
 - Solid state drives are emerging for smaller scratch space
 - Tape systems for archiving

Login nodes			
ᆕ		Storage	;
	Interconne	ect 🤇	
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 - $\circ \ \ \, {\rm Still} \ \, {\rm predominantly} \ \, {\rm spinning} \ \, {\rm disks}$
 - Solid state drives are emerging for smaller scratch space
 - Tape systems for archiving
- Front-end nodes
 - For user access
 - Compiling, submitting jobs, etc.



Specific configuration of the Cyl cluster

- Various nodes with different architectures
 - Hostnames: cyc0{1, ...,8}, ph0{1,2,3,4}, cwp0{1,2}, cwg0{1,2}, sim0?, etc.
- We will be using nodes from the p100 partition for now:
 - Hostnames are cyc01, cyc02, ..., cyc08. Three are reserved for our lesson
 - \circ 2×16-core Intel Xeon
 - 128 GBytes RAM
 - $\circ~2{\times}P100\,GPUs$ each
- Common storage for our course: /onyx/data/sds406f24/



- Log in to a login node or frontend node. Login node in our case has hostname front02
- To run programs on *compute nodes*, a *job scheduler* is available
- Distinguish between *interactive* and *batch* jobs

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SLURM job scheduler

- See currently running and waiting jobs: squeue
- Ask for an interactive job: salloc
- Submit a batch job: sbatch
- Run an executable: srun

• Log in:

[localhost ~]\$ ssh <username>@front02.hpcf.cyi.ac.cy

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• Type hostename. This tells you the name of the node you are currently logged into:

[ikoutsou@front02 ~]\$ hostname
front02

this is the login node.

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• Type hostename. This tells you the name of the node you are currently logged into:

[ikoutsou@front02 ~]\$ hostname
front02

this is the login node.

• Ask for a node:

```
[ikoutsou@front02 ~]$ salloc -N 1 -p p100 --reservation=sds406 -A sds406f24
salloc: Granted job allocation 298614
[ikoutsou@cyc01 ~]$
```

• Type hostname again:

[ikoutsou@cyc01 ~]\$ hostname
cyc01

this is a compute node.

• Release the node:

[ikoutsou@cyc01 ~]\$ exit
exit
salloc: Relinquishing job allocation 69053
[ikoutsou@front02 ~]\$

we're back on front02

• Release the node:

[ikoutsou@cyc01 ~]\$ exit
exit
salloc: Relinquishing job allocation 69053
[ikoutsou@front02 ~]\$

we're back on front02

• Our course reservation includes the nodes with cyc in the hostname.

• Release the node:

[ikoutsou@cyc01 ~]\$ exit
exit
salloc: Relinquishing job allocation 69053
[ikoutsou@front02 ~]\$

we're back on front02

• Our course reservation includes the nodes with cyc in the hostname.

Please do not hold nodes unnecessarily; when you have nodes salloced you may be blocking other users from using those nodes.

• Release the node:

[ikoutsou@cyc01 ~]\$ exit
exit
salloc: Relinquishing job allocation 69053
[ikoutsou@front02 ~]\$

we're back on front02

• Our course reservation includes the nodes with cyc in the hostname.

Please do not hold nodes unnecessarily; when you have nodes salloced you may be blocking other users from using those nodes.

• Use srun instead of salloc:

```
[ikoutsou@front02 ~]$ srun -n 1 -N 1 -p p100 --reservation=sds406 -A sds406f24 hostname
srun: job 203373 queued and waiting for resources
srun: job 203373 has been allocated resources
cyc01
```

Allocates a node, runs the specified command (in this case hostname), and then exits the node, releasing the allocation.

• Run multiple instances of hostname in parallel:

[ikoutsou@front02 ~]\$ srun -N 1 -n 2 -p p100 --reservation=sds406 -A sds406f24 hostname
srun: job 203374 queued and waiting for resources
srun: job 203374 has been allocated resources
cyc01
cyc01

• Run multiple instances of hostname in parallel:

[ikoutsou@front02 ~]\$ srun -N 1 -n 2 -p p100 --reservation=sds406 -A sds406f24 hostname
srun: job 203374 queued and waiting for resources
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- \circ -N 1: use one node
- -n 2: use two processes

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srun: job 203374 has been allocated resources
cyc01
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- \circ -N 1: use one node
- -n 2: use two processes
- Run on more than one node:

```
[ikoutsou@front02 ~]$ srun -N 2 -n 2 -p p100 --reservation=sds406 -A sds406f24 hostname
srun: job 203375 queued and waiting for resources
srun: job 203375 has been allocated resources
cyc01
cyc02
```

runs one instance of hostname on each node

• Run multiple instances of hostname in parallel:

```
[ikoutsou@front02 ~]$ srun -N 1 -n 2 -p p100 --reservation=sds406 -A sds406f24 hostname
srun: job 203374 queued and waiting for resources
srun: job 203374 has been allocated resources
cyc01
cyc01
```

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srun: job 203375 queued and waiting for resources
srun: job 203375 has been allocated resources
cyc01
cyc02
```

runs one instance of hostname on each node

• Try:

[ikoutsou@front02 ~]\$ srun -N 2 -n 4 -p p100 --reservation=sds406 -A sds406f24 hostname

[ikoutsou@front02 ~]\$ srun -N 2 -n 3 -p p100 --reservation=sds406 -A sds406f24 hostname

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- Emacs and Vim are available on the cluster
- Other options are fine, as long as you know what you're doing

• Make a directory. List it.

[ikoutsou@front02 ~]\$ mkdir SDS406 [ikoutsou@front02 ~]\$ ls SDS406 [ikoutsou@front02 ~]\$

• Make a directory. List it.

```
[ikoutsou@front02 ~]$ mkdir SDS406
[ikoutsou@front02 ~]$ ls
SDS406
[ikoutsou@front02 ~]$
```

• Change into it:

[ikoutsou@front02 ~]\$ cd SDS406/ [ikoutsou@front02 SDS406]\$

• Make a directory. List it.

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[ikoutsou@front02 ~]$ mkdir SDS406
[ikoutsou@front02 ~]$ ls
SDS406
[ikoutsou@front02 ~]$
```

• Change into it:

```
[ikoutsou@front02 ~]$ cd SDS406/
[ikoutsou@front02 SDS406]$
```

• While in SDS406/, make another one for this week's lesson and change into it:

[ikoutsou@front02 SDS406]\$ mkdir l01 [ikoutsou@front02 SDS406]\$ cd l01/ [ikoutsou@front02 l01]\$

• Make a directory. List it.

```
[ikoutsou@front02 ~]$ mkdir SDS406
[ikoutsou@front02 ~]$ ls
SDS406
[ikoutsou@front02 ~]$
```

• Change into it:

```
[ikoutsou@front02 ~]$ cd SDS406/
[ikoutsou@front02 SDS406]$
```

• While in SDS406/, make another one for this week's lesson and change into it:

[ikoutsou@front02 SDS406]\$ mkdir l01
[ikoutsou@front02 SDS406]\$ cd l01/
[ikoutsou@front02 l01]\$

• pwd will tell you where you are in the file system:

[ikoutsou@front02 l01]\$ pwd
/home/ikoutsou/SDS406/l01

Let's write our own hostname command (you were hoping for "Hello world" 🙃 ?)

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[ikoutsou@front02 l01]\$ emacs my_hn.c

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• Type in the following program:

```
#include <unistd.h>
#include <stdio.h>
int
main(int argc, char *argv[])
{
     char hname[256];
     gethostname(hname, 256);
     printf(" Hostname is: %s\n", hname);
     return 0;
}
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- To save, hold down ctrl, then hold and release x, then hold and release s
- To exit emacs, hold down ctrl, then hold and release x, then hold and release c

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- To exit emacs, hold down ctrl, then hold and release x, then hold and release c
- Bookmark the Emacs reference card: <u>https://www.gnu.org/software/emacs/refcards/pdf/refcard.pdf</u>

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• 1s should show the file my_hn.c, which you just typed in and saved:

[ikoutsou@front02 l01]\$ ls my_hn.c

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• It's time to *compile* it into an executable we can run:

[ikoutsou@front02 l01]\$ module load gompi
[ikoutsou@front02 l01]\$ gcc my_hn.c -o hn

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• -• hn means "name the resulting executable hn". If you don't specify -• the executable name defaults to a.out

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• Type ls to make sure it has been created. Then run it on the frontend node:

[ikoutsou@front02 l01]\$ ls
hn my_hn.c
[ikoutsou@front02 l01]\$./hn
Hostname is: front02
Cluster Computing Introductory Example

After exiting emacs you are back at the command line

• 1s should show the file my_hn.c, which you just typed in and saved:

[ikoutsou@front02 l01]\$ ls
my_hn.c

• It's time to *compile* it into an executable we can run:

```
[ikoutsou@front02 l01]$ module load gompi
[ikoutsou@front02 l01]$ gcc my_hn.c -o hn
```

• -• hn means "name the resulting executable hn". If you don't specify -• the executable name defaults to a.out

• Type ls to make sure it has been created. Then run it on the frontend node:

```
[ikoutsou@front02 l01]$ ls
hn my_hn.c
[ikoutsou@front02 l01]$ ./hn
Hostname is: front02
```

Note: commands like gcc or ls are globally accessible because their locations are included in your shell environment's search path. For hn though, which we just created, you need to explicitly give its path, in this case via . / which means "current directory".