Introduction to High Performance Computing

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SDS406 – Fall semester, 2024 - 2025

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L08: Distributed memory parallel programming, 18th November 2024

Outline

Parallel programming

- Distributed memory paradigm
- Main features and relation to shared memory parallelism

MPI Introduction

- The MPI standard and development workflow
- Using MPI
- Hands-on examples

The Message Passing Interface

- MPI: An Application Programmer Interface (API)
 - A library specification; determines functions, their names and arguments, and their functionality
- A *de facto* standard for programming *distributed memory* systems
- Current specification is version 4 (MPI-4.0), released June 9, 2021
 - $\circ~$ For most systems you can reliably assume MPI-3.1 is in place
- Several free (open) or vendor-provided implementations, e.g.:
 - Mvapich
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Distributed memory programming

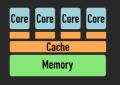
- Each process has its own memory domain
- MPI functions facilitate:
 - Obtaining environment information about the running process, e.g., process id, number of processes, etc.
 - Achieving *communication* between processes, e.g. synchronization, copying of data, etc.

Shared memory



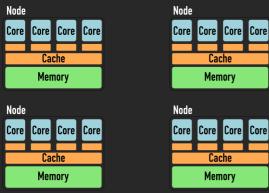
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- E.g. multi-core CPU, multi-socket node, GPU threads, etc.
- Programming models: OpenMP, pthreads, MPI, CUDA (sort of)

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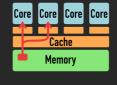
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- E.g. multiple nodes within a cluster, multiple GPUs within a node
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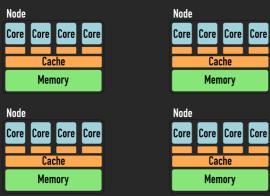
Shared memory



Data shared via memory

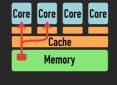
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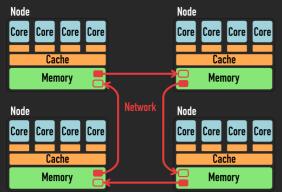
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Distributed memory



Data shared via explicit communication over a network

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- E.g. multiple nodes within a cluster, multiple GPUs within a node
- Programming models: MPI

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• Depending on the system, instead of mpirun you may be required mpiexec or srun which take similar (but not identical) arguments

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• For Python, a Python module, mpi4py, implements the MPI API

import mpi4py

or, more commonly

from mpi4py import MPI

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#include <mpi.h>
int
main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);
    /*
    ...
    ...
    ...
    */
    MPI_Finalize();
    return 0;
}
```

- Two functions you will almost always call
 - MPI_Comm_size() or MPI.COMM_WORLD.Get_size(): gives the number of parallel process running (n_{proc})
 - MPI_Comm_rank() or MPI.COMM_WORLD.Get_rank(): determines the rank of the process, i.e. a unique number between 0 and $n_{proc} 1$ that identifies the calling process
- A complete example:

```
#include <stdio.h>
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int
main(int argc, char *argv[])
{
    MPI_Init(&argc, &argv);
    int nproc, rank;
    MPI_Comm_size(MPI_COMM_WORLD, &nproc);
    MPI_Comm_rank(MPI_COMM_WORLD, &rank);
    printf(" This is rank = %d of nproc = %d\n", rank, nproc);
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- MPI_COMM_WORLD or MPI.COMM_WORLD is an MPI communicator
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 - By default, the initial communicator is COMM_WORLD, i.e. all processes in one communicator

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 - By default, the initial communicator is COMM_WORLD, i.e. all processes in one communicator
- No assumptions can safely be made about the order in which the printf() statements occur, i.e. the order in which each process prints is practically random

• Compiling and running the previous program (assuming it is saved as example.c or example.py)

```
[user@front02 ~]$ mpicc -o example example.c
[user@front02 ~]$ mpirun -n 5 example
This is rank = 3 of nproc = 5
This is rank = 1 of nproc = 5
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This is rank = 4 of nproc = 5
This is rank = 0 of nproc = 5
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[user@front02 ~]$ mpirun -n 5 python example.py
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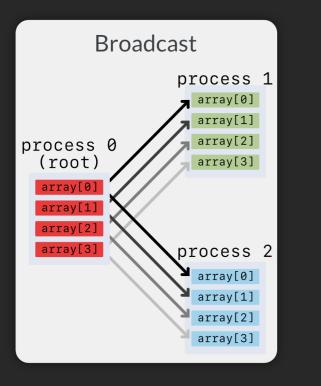
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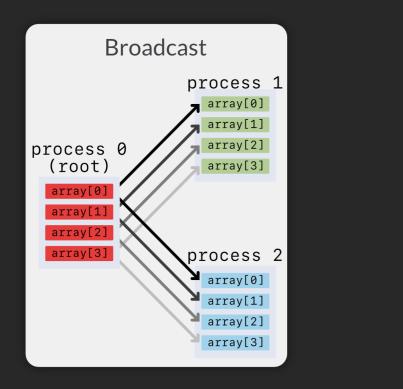
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- Most collective operations implicitly synchronize the processes

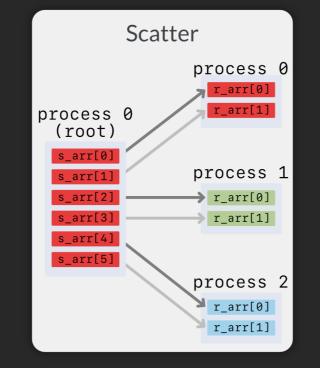
- The first set of communication functions we will look at are collective operations
- Collective: all processes must be involved in the operation (as opposed to point-to-point communications)
- Examples (this list is not exhaustive!):
 - Broadcast a variable from one process to all processes (Broadcast)
 - Distribute elements of an array on one process to multiple processes (Scatter)
 - Collect elements of arrays scattered over processes into a single process (Gather)
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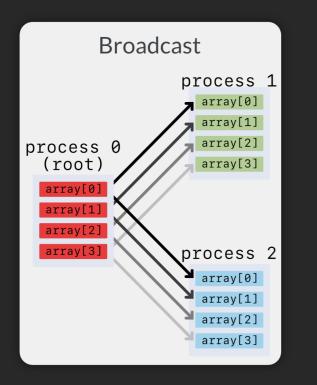


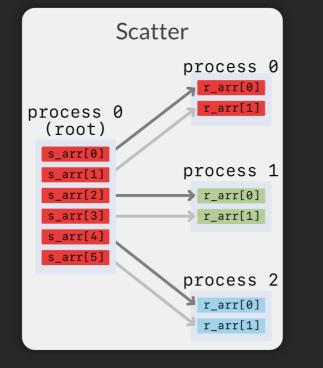
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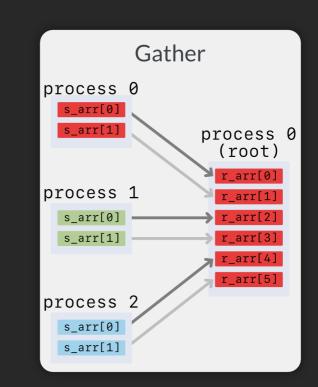




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- In C, we can pass a scalar by using its memory address (second example above). In Python, we can use zero-dimensional numpy array for this
- Full list of types available in MPI documentation. E.g. see: <u>https://www.mpich.org/static/docs/latest/www3/Constants.html</u>

• Scatter:

MPI_Scatter(
 const void *sendbuf, int sendcount, MPI_Datatype sendtype,
 void *recvbuf, int recvcount, MPI_Datatype recvtype,
 int root, MPI_Comm comm
);

MPI.COMM_WORLD.Scatter(

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- Example: distribute a 12-element array from process 0, assuming 3 processes in total (including root)

```
arr_all = np.random.rand(12)
arr_proc = np.zeros([4])
MPI.COMM WORLD.Scatter(arr all, arr proc, root = 0)
```

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    int root, MPI_Comm comm
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```

• Example: distribute each element of a 4-element array to 4 processes in total (including root)

```
arr = np.random.rand(4)
element = np.zeros([])
MPI.COMM_WORLD.Scatter(arr, element, root = 0)
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Note: the initialization of element as a zero-dimensional array

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• Gather:

• Example: collect a 9-element array at process 0, by concatenating 3 elements from each of 3 processes in total (including root)

```
arr_all = np.zeros([9])
arr_proc = np.random.rand(3)
MPI.COMM_WORLD.Gather(arr_proc, arr_all, root = 0)
```

• Gather:

• Example: collect a 4-element array at process 0, by concatenating an element from each of 4 processes in total (including root)

```
double arr[4]; /* ← this only needs to be defined on process with rank = 0 */
double element;
MPI_Gather(&element, 1, MPI_DOUBLE, arr, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);
```

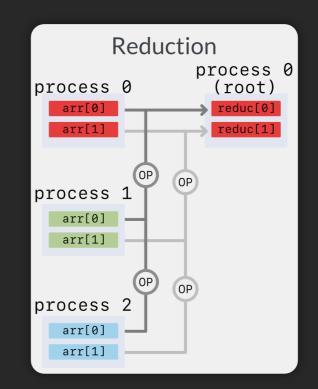
```
arr = np.zeros([4])
element = np.random.rand(1)
MPI.COMM_WORLD.Gather(element, arr, root = 0)
```

• Reduction:

MPI_Reduce(

const void *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm MPI.COMM_WORLD.Reduce(

sendbuf: BufSpec, recvbuf: BufSpec,
op: Op = SUM, root: int = 0



• Reduction:

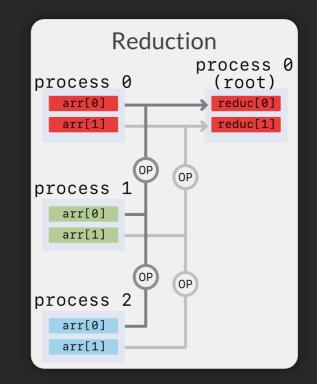
MPI_Reduce(

const void *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm
)

• Notes:

- Op is an operation, e.g. MPI_SUM, MPI_PROD, etc. (MPI.SUM and MPI.PROD in Python)
- In C note the need for specifying the datatype and count, the number of elements of the arrays
- The operation is over all processes in comm, in this case COMM_WORLD

MPI.COMM_WORLD.Reduce(



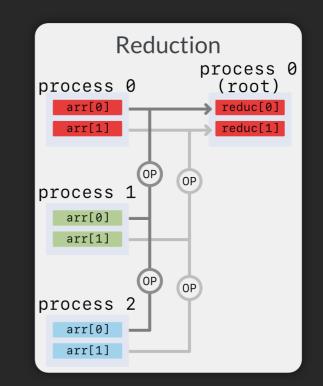
• Reduction:

MPI_Reduce(

const void *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm
)

• *Example*: Sum each element of a 3-element array over all processes

```
s_arr = np.random.rand(3)
r_arr = np.zeros([3])
MPI.COMM_WORLD.Reduce(s_arr, r_arr, MPI.SUM, root = 0)
```



• Reduction:

MPI_Reduce(

const void *sendbuf, void *recvbuf, int count, MPI_Datatype datatype, MPI_Op op, int root, MPI_Comm comm
)

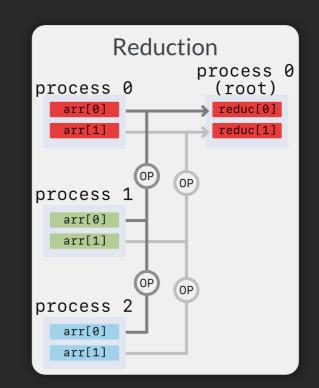
• *Example*: Sum variable var over all processes

double var; double sum; /* ← only needs to * * be defined on root */ MPI_Reduce(&var, &sum, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD);

var = np.random.rand(1)
sum = np.zeros([])
MPI.COMM_WORLD.Reduce(var, sum, MPI.SUM, root = 0)

MPI.COMM_WORLD.Reduce(

sendbuf: BufSpec, recvbuf: BufSpec,
op: Op = SUM, root: int = 0



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 - E.g. below, the sum will be placed in var of the root process (process with rank = 0):

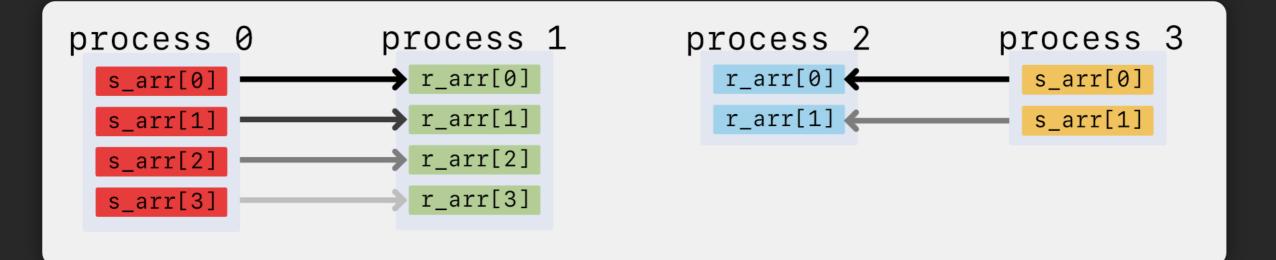
```
if(rank ≠ 0) {
    MPI_Reduce(&var, null, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD);
} else {
    MPI_Reduce(MPI_IN_PLACE, &var, 1, MPI_DOUBLE, MPI_SUM, 0, MPI_COMM_WORLD);
}
```

• Communications that involve transfer of data between two processes

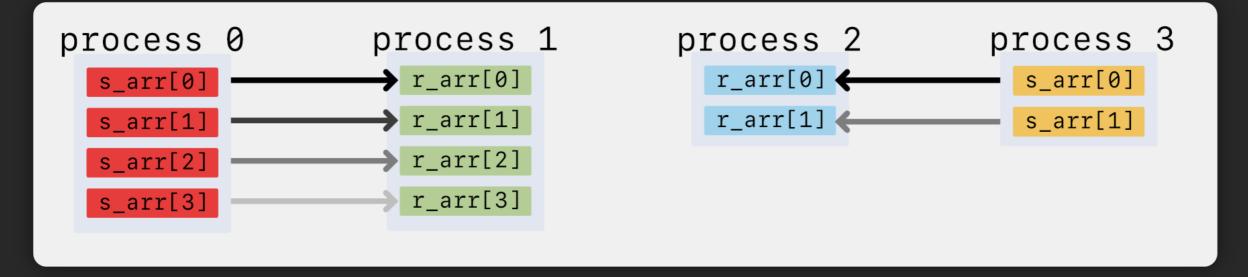
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Two point-to-point communications are depicted above between i) process 0 and 1 and between ii) process 2 and 3

• Send/Receive

MPI_Send(void *buf, int n, MPI_Datatype type, int dest, int tag, MPI_Comm comm)
MPI_Recv(void *buf, int n, MPI_Datatype type, int srce, int tag, MPI_Comm comm, MPI_Status *status)

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MPI.COMM_WORLD.Send(buf: BufSpec, dest: int, tag: int = 0)
MPI.COMM_WORLD.Recv(buf: BufSpec, source: int = ANY_SOURCE, tag: int = ANY_TAG, status: Status = None)

• Note the need to specify a source and destination rank

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- Use of MPI_ANY_SOURCE (MPI.ANY_SOURCE in Python) in Recv() means "accept data from any source"
- status can be used to query the result of the receive (e.g. how many elements were received). We will set to MPI_STATUS_IGNORE

• Send/Receive; a trivial example



• Send/Receive; a trivial example



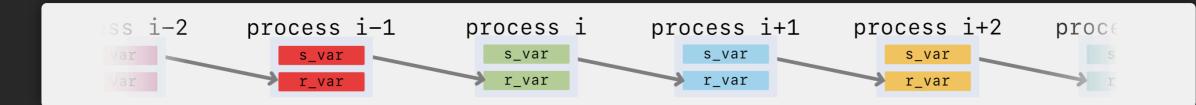
```
if(rank = i) {
    MPI_Send(s_arr, 4, MPI_DOUBLE, j, 0, MPI_COMM_WORLD);
}
if(rank = j) {
    MPI_Recv(r_arr, 4, MPI_DOUBLE, i, MPI_ANY_TAG, MPI_COMM_WORLD, MPI_S
}
```

```
s_arr = np.random.rand(4)
r_arr = np.zeros([4])
if rank = i:
    MPI.COMM_WORLD.Send(s_arr, j)
if rank = j:
    MPI.COMM_WORLD.Recv(r_arr, i)
```

• It is common in parallel applications to require that every process communicates with another process, e.g. a neighboring process



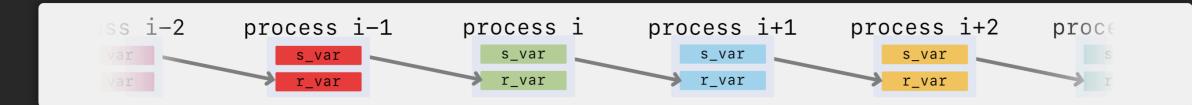
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• This will **not** work:

MPI_Send(&s_var, 1, MPI_DOUBLE, rank+1, 0, MPI_COMM_WORLD); MPI_Recv(&r_var, 1, MPI_DOUBLE, rank-1, MPI_ANY_TAG, MPI_COMM_WORLD, M MPI.COMM_WORLD.Send(s_var, rank+1); MPI.COMM_WORLD.Recv(r_var, rank-1);

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MPI.COMM WORLD.Send(s var, rank+1);

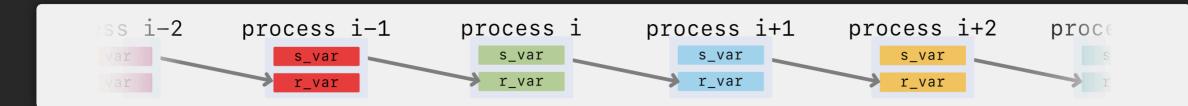
MPI.COMM WORLD.Recv(r var, rank-1);

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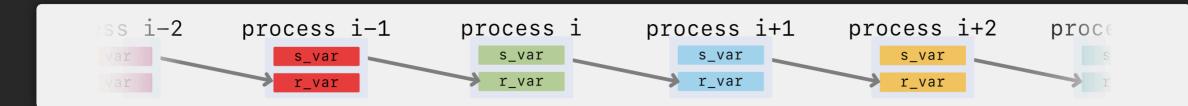
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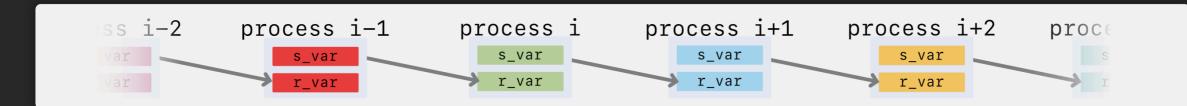
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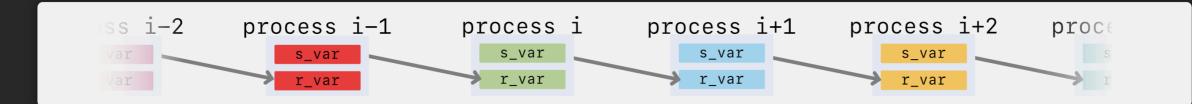
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 - Serializes communications that would otherwise be done faster in parallel
 - Inelegant, obscure, and error-prone

MPI.COMM_WORLD.Send(s_var, rank+1); MPI.COMM WORLD.Recv(r var, rank-1);

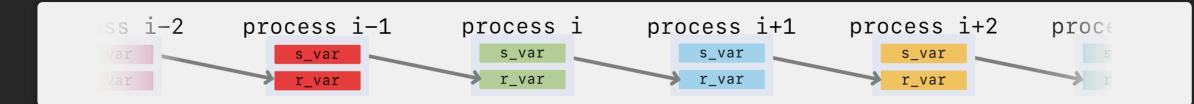
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• A more efficient and elegant solution is to use Sendrecv():

```
MPI.COMM_WORLD.Sendrecv(
    sendbuf: BufSpec, dest: int, sendtag: int = 0,
    recvbuf: BufSpec, source: int = ANY_SOURCE, recvtag: int = ANY_TAG,
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• For the depicted example:

MPI_Sendrecv(&s_var, 1, MPI_DOUBLE, rank+1, 0, &r_var, 1, MPI_DOUBLE, rank-1, MPI_ANY_TAG, MPI_COMM_WORLD, MPI_STATUS_IGNORE);

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MPI_Isend(sendbuf, ..., request);
    /*
    * More code can come here, provided it
    * does not modify sendbuf, which is
    * assumed to be "in-flight"
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MPI_Wait(request, ...);
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Some additional notes on variants of the point-to-point communications we have covered

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```

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• Like MPI_Sendrecv() but with a single buf rather than separate sendbuf and recvbuf The receive message overwrites the send message

[•] MPI_Sendrecv_replace()

cp -r /onyx/data/sds406f24/l08/ex?? .

• Exercises follow a common structure

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 - ex04: Use of Bcast(), Scatter(), Sendrecv(), and Gather()
- All exercises have been tested with specific versions of OpenMPI and the GNU Compiler. Please use:

module load gompi/2023a

for all exercises.

Ex01

• Modify ex01.c to call MPI_Comm_size() and MPI_Comm_rank() appropriately

/*
 * TODO: call the appropriate MPI functions here
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• A job script has been prepared to run ex01:

[user@front02 ex01]\$ cat sub-ex01.sh #!/bin/bash #SBATCH -- job-name=01 #SBATCH -- nodes=2 #SBATCH -- ntasks=8 #SBATCH -- ntasks-per-node=4 #SBATCH -- output=ex01-output.txt #SBATCH -- time=00:01:00 #SBATCH -- partition=p100 #SBATCH -A sds406f24 module load gompi/2023a

mpirun ex01

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• Modify ex01.c to call MPI_Comm_size() and MPI_Comm_rank() appropriately

```
/*
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 */
```

Compile

mpicc -o ex01 ex01.c

• A job script has been prepared to run ex01:

```
[user@front02 ex01]$ cat sub-ex01.sh
#!/bin/bash
#SBATCH -- job-name=01
#SBATCH -- nodes=2
#SBATCH -- ntasks=8
#SBATCH -- ntasks-per-node=4
#SBATCH -- output=ex01-output.txt
#SBATCH -- time=00:01:00
#SBATCH -- partition=p100
#SBATCH -A sds406f24
module load gompi/2023a
```

mpirun ex01

• 2 nodes, 8 processes, meaning 4 processes per node

Ex01

• Modify ex01.c to call MPI_Comm_size() and MPI_Comm_rank() appropriately

```
/*

* TODO: call the appropriate MPI functions here

*/
```

Compile

mpicc -o ex01 ex01.c

• A job script has been prepared to run ex01:

```
[user@front02 ex01]$ cat sub-ex01.sh
#!/bin/bash
#SBATCH -- job-name=01
#SBATCH -- nodes=2
#SBATCH -- ntasks=8
#SBATCH -- ntasks-per-node=4
#SBATCH -- output=ex01-output.txt
#SBATCH -- time=00:01:00
#SBATCH -- partition=p100
#SBATCH -A sds406f24
madule_lead_gempi(2022a)
```

module load gompi/2023a mpirun ex01

- 2 nodes, 8 processes, meaning 4 processes per node
- program output will be redirected to file ex01-output.txt

Ex01

• Submit the job script:

[user@front02 ex01]\$ sbatch sub-ex01.sh Submitted batch job 69711 [user@front02 ex01]\$

Ex01

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Ex01

• Submit the job script:

[user@front02 ex01]\$ sbatch sub-ex01.sh Submitted batch job 69711 [user@front02 ex01]\$

- If done, the file ex01-output.txt should have been created
- Inspect the file:

[user@front02 ex01]\$ cat ex01-output.txt This is rank = 0 of nproc = 8 on node: cyc06 This is rank = 2 of nproc = 8 on node: cyc06 This is rank = 3 of nproc = 8 on node: cyc06 This is rank = 5 of nproc = 8 on node: cyc07 This is rank = 6 of nproc = 8 on node: cyc07 This is rank = 1 of nproc = 8 on node: cyc06 This is rank = 7 of nproc = 8 on node: cyc07 This is rank = 4 of nproc = 8 on node: cyc07 [user@front02 ex01]\$

Ex01

• Submit the job script:

[user@front02 ex01]\$ sbatch sub-ex01.sh Submitted batch job 69711 [user@front02 ex01]\$

- If done, the file ex01-output.txt should have been created
- Inspect the file:

[user@front02 ex01]\$ cat ex01-output.txt This is rank = 0 of nproc = 8 on node: cyc06 This is rank = 2 of nproc = 8 on node: cyc06 This is rank = 3 of nproc = 8 on node: cyc06 This is rank = 5 of nproc = 8 on node: cyc07 This is rank = 6 of nproc = 8 on node: cyc07 This is rank = 1 of nproc = 8 on node: cyc06 This is rank = 7 of nproc = 8 on node: cyc07 This is rank = 4 of nproc = 8 on node: cyc07 [user@front02 ex01]\$

• Note the order is nondeterministic; whichever process reaches the print statement first prints

Ex02

- ex02.py demonstrates the use of Scatter() and Reduce()
- The file with name array.txt includes 55,440 floating point numbers, one per line:

[user@front02 ex02]\$ head array.txt 7.913676052329088328e-01 1.879167007836126668e-01 2.343674804515035737e-01 4.707043244181141617e-02 6.272795840838938375e-01 2.725799268304553991e-01 5.803516013116442052e-01 2.356271465482765448e-01 2.982738904468156260e-01 5.372364132030218453e-01 [ikoutsou@front02 ex02]\$

Ex02

• We would like:

Ex02

• We would like:

• The root process to read all elements into an array array[]

Ex02

• We would like:

- The root process to read all elements into an array array[]
- The root process to *broadcast* the total number of elements of the array, ntot

Ex02

• We would like:

- The root process to read all elements into an array array[]
- The root process to *broadcast* the total number of elements of the array, ntot
- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size

Ex02

- The root process to read all elements into an array array[]
- The root process to *broadcast* the total number of elements of the array, ntot
- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size
- The root process to scatter the elements of array array[] to all processes
 - └→ Each process should receive nloc elements

Ex02

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- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size
- The root process to scatter the elements of array array[] to all processes
 Each process should receive nloc elements
- Each process to sum its local elements, storing the result into <code>sum_loc</code>

Ex02

- The root process to read all elements into an array <code>array[]</code>
- The root process to *broadcast* the total number of elements of the array, ntot
- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size
- The root process to scatter the elements of array array[] to all processes
 Each process should receive nloc elements
- $\circ~$ Each process to sum its local elements, storing the result into ${\tt sum_loc}$
- $\circ~$ To use a reduction operation to obtain the grand total over all 55,440 elements in the root rank

Ex02

- The root process to read all elements into an array array[]
- The root process to *broadcast* the total number of elements of the array, ntot
- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size
- The root process to scatter the elements of array array[] to all processes
 Each process should receive nloc elements
- $\circ~$ Each process to sum its local elements, storing the result into <code>sum_loc</code>
- To use a reduction operation to obtain the grand total over all 55,440 elements in the root rank
- Look at ex02.py. You only need to complete some parts, as instructed by the comments with TODO

Ex02

• We would like:

- The root process to read all elements into an array array[]
- The root process to broadcast the total number of elements of the array, ntot
- Each process to initialize an empty array sub[] with number of elements nloc = ntot / size
- The root process to scatter the elements of array array[] to all processes
 Each process should receive nloc elements
- Each process to sum its local elements, storing the result into <code>sum_loc</code>
- To use a reduction operation to obtain the grand total over all 55,440 elements in the root rank
- Look at ex02.py. You only need to complete *some* parts, as instructed by the comments with TODO
- The correct result, which will be in ex02-output.txt should be:

Sum: 27777.25711

Ex03

- This exercise demonstrates Gather()
- A file filenames.txt includes the filenames of 8 files:

[user@front02 ex03]\$ cat filenames.txt 00.txt 01.txt 02.txt 03.txt 04.txt 05.txt 06.txt 07.txt [user@front02 ex03]\$

Ex03

- This exercise demonstrates Gather()
- A file filenames.txt includes the filenames of 8 files:

[user@front02 ex03]\$ cat filenames.txt 00.txt 01.txt 02.txt 03.txt 04.txt 05.txt 06.txt 07.txt [user@front02 ex03]\$

• In ex03.c, the root process (process with rank = 0) reads the filenames and scatters one to each process

Ex03

- This exercise demonstrates Gather()
- A file filenames.txt includes the filenames of 8 files:

[user@front02 ex03]\$ cat filenames.txt 00.txt 01.txt 02.txt 03.txt 04.txt 05.txt 06.txt 07.txt [user@front02 ex03]\$

- In ex03.c, the root process (process with rank = 0) reads the filenames and scatters one to each process
- Each process then computes the Fletcher 32 checksum of one file

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- A file filenames.txt includes the filenames of 8 files:

[user@front02 ex03]\$ cat filenames.txt 00.txt 01.txt 02.txt 03.txt 04.txt 05.txt 06.txt 07.txt [user@front02 ex03]\$

- In ex03.c, the root process (process with rank = 0) reads the filenames and scatters one to each process
- Each process then computes the Fletcher 32 checksum of one file
- You need to write an appropriate Gather() operation to collect the checksums into the root process such that it prints them correctly

Ex03

• See also ex03-scalar.c, which implements the same but with no parallelism,

[user@front02 ex03]\$ srun -n 1 -p p100 ./ex03-scalar filenames.txt * 00.txt → 04D70552 * 01.txt → 19708CD4 * 02.txt → ED737A1C * 03.txt → 0C40E2D2 * 04.txt → F7BDE74D * 05.txt → 562DDD6C * 06.txt → 6F2CD2F1 * 07.txt → 016DB6C6 Done 8 files in t = 12.953475 sec [user@front02 ex03]\$

Ex04

• In this exercise, a large array of 17,463,600 elements is read by rank 0 (n_arr = 17463600)

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- Our objective is to compute a second order discrete derivative of the data:

deriv[i] = arr[i-1] + arr[i+1] - 2*arr[i]

• The goal is to do this in parallel, such that each rank computed it for a subset of the data

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- We will proceed as follows:

Ex04

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 - The array in each rank has two additional elements, arr_loc[n_loc+2], where n_loc = n_arr / nproc

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- We will proceed as follows:
 - The array in each rank has two additional elements, arr_loc[n_loc+2], where n_loc = n_arr / nproc
 - Add a missing Scatter() to distribute to each processes the corresponding n_loc elements. These should be received in arr_loc[1:-1] of each rank

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 - Now your program can proceed to correctly compute the derivative:

deriv_loc[i] = arr_loc[i-1] - 2*arr_loc[i] + arr_loc[i+1]

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• Use a Gather() to collect the full array of the derivative into the root process (rank = 0). The root process will then write the array to a new file deriv.txt

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• If done correctly, deriv.txt will contain all +2.000000, except for the first and last element

	rank	0
array.txt	array[::]
0.000000000	array[
0.0000000000	array[1
2.0000000000	array[1 2 3
6.0000000000	array[
12.000000000	array[4
20.000000000	array[5
•	•	
•	•	
•	•	
•		
•	array[-3
•	array[
•	array[-1.

